Modal proof theory through a focused telescope

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Proof theory

Mathematicians consider mathematical objects (groups, vector spaces, ...)

Proof:

1. a convincing argument, a token of evidence, ...

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Proof theorists consider proofs as their objects

Object proofs

Proof:

- 1. a convincing argument, a token of evidence, \ldots
- 2. a list of inferences starting with known facts (*axioms*), combining them, and ending on a new conclusion (*theorem*)

Proof theorists consider proofs as their objects and prove some properties about them.

Object proofs

Meta-proofs





Natalie Dee.com













Outline

CARefully Illustrating NEsted sequents

Logic Modal logic Modal proof theory

A Unique REarrangement of LIsted ENtities

Proof search Focusing Synthetic rules

Actual New Utterances (with Passion And Method)

Folding Unfolding

CARefully Illustrating NEsted sequents

Propositional logic:

$$p$$
 : it is raining $\widehat{\mathbb{M}}_{atomic \ propositions} \left\{ ar{p} : \text{ it is not raining } \mathfrak{S} \ s \ not \ sad \ \odot \right\}_{atomic \ propositions} \left\{ ar{s} : she \ is \ not \ sad \ \odot \right\}$

Propositional logic:

$$\begin{array}{rcl} p & : & \text{it is raining } \widehat{\mathbb{W}} \\ s & : & \text{she is sad } \odot \end{array} \right\}_{\text{atomic propositions}} \left\{ \begin{array}{c} \overline{p} & : & \text{it is not raining } \overset{\mathfrak{G}}{\mathbb{S}} \\ \overline{s} & : & \text{she is not sad } \odot \end{array} \right\}$$

 $p \supset s$: if it is raining then she is sad

Propositional logic:

p s	:	it is raining $\widehat{\mathbb{M}}$ at she is sad \mathbb{S}	tomic propositions $\begin{cases} ar{p} \\ ar{s} \end{cases}$:	it is not raining 🖑 she is not sad ©

- $p \supset s$: if it is raining then she is sad
- $p \land s$: it is raining and she is sad

Propositional logic:

 $\begin{array}{rcl} p & : & it is raining & \\ s & : & she is sad & \\ \end{array}\right\}_{atomic \ propositions} \left\{ \begin{array}{rcl} \overline{p} & : & it is not raining & \\ \hline{s} & : & she is not sad & \\ \end{array} \right\}$ $p \supset s & : & if it is raining then she is sad \\ p \land s & : & it is raining and she is sad \\ \end{array}$

 $p \lor s$: it is raining or she is sad

Propositional logic:

 $\begin{array}{rcl} p & : & it \ is \ raining & \\ s & : & she \ is \ sad & \\ \hline p \supset s & : & if \ it \ is \ raining \ then \ she \ is \ sad \\ p \land s & : & it \ is \ raining \ and \ she \ is \ sad \\ p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right\}^{atomic \ propositions} \left\{ \begin{array}{rcl} \overline{p} & : & it \ is \ not \ raining & \\ \hline \overline{s} & : & she \ is \ not \ sad & \\ \hline s & : & she \ is \ not \ sad & \\ \hline p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right.$

Modal logic:

Propositional logic:

 $\begin{array}{rcl} p & : & it \ is \ raining & \\ s & : & she \ is \ sad & \\ \hline p \supset s & : & if \ it \ is \ raining \ then \ she \ is \ sad \\ p \land s & : & it \ is \ raining \ and \ she \ is \ sad \\ p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right\}^{atomic \ propositions} \left\{ \begin{array}{c} \overline{p} & : & it \ is \ not \ raining & \\ \hline \overline{s} & : & she \ is \ not \ sad & \\ \hline s & : & she \ is \ not \ sad & \\ \hline p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right.$

Modal logic:

 $\Diamond p$: it is possible that it is raining

Propositional logic:

 $\begin{array}{rcl} p & : & it \ is \ raining & \\ s & : & she \ is \ sad & \\ \hline p \supset s & : & if \ it \ is \ raining \ then \ she \ is \ sad \\ p \land s & : & it \ is \ raining \ and \ she \ is \ sad \\ p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right\}^{atomic \ propositions} \left\{ \begin{array}{c} \overline{p} & : & it \ is \ not \ raining & \\ \hline \overline{s} & : & she \ is \ not \ sad & \\ \hline s & : & she \ is \ not \ sad & \\ \hline p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right.$

Modal logic:

- $\Diamond p$: it is possible that it is raining
- $\Box s$: it is necessary that she is sad

Propositional logic:

 $\begin{array}{rcl} p & : & it \ is \ raining & \\ s & : & she \ is \ sad & \\ \hline p \supset s & : & if \ it \ is \ raining \ then \ she \ is \ sad \\ p \land s & : & it \ is \ raining \ and \ she \ is \ sad \\ p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right\}^{atomic \ propositions} \left\{ \begin{array}{rcl} \overline{p} & : & it \ is \ not \ raining & \\ \hline \overline{s} & : & she \ is \ not \ sad & \\ \hline s & : & she \ is \ not \ sad & \\ \hline p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right.$

Modal logic:

¢р	:	it is possible that it is raining
5	:	it is necessary that she is sad

First-order logic:

I(x)	:	she loves x	
h(x, y)	:	x hates y	atomic predicates

Propositional logic:

 $\begin{array}{rcl} p & : & it \ is \ raining & \\ s & : & she \ is \ sad & \\ \hline p \supset s & : & if \ it \ is \ raining \ then \ she \ is \ sad \\ p \land s & : & it \ is \ raining \ and \ she \ is \ sad \\ p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right\}^{atomic \ propositions} \left\{ \begin{array}{rcl} \overline{p} & : & it \ is \ not \ raining & \\ \hline \overline{s} & : & she \ is \ not \ sad & \\ \hline s & : & she \ is \ not \ sad & \\ \hline p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right.$

Modal logic:

¢р	:	it is possible that it is raining
<u>s</u>	:	it is necessary that she is sad

First-order logic:

l(x) h(x,y)	:	$\left. \begin{array}{c} she \ loves \ x \\ x \ hates \ y \end{array} \right\}_{\text{atomic predicates}}$
∀ x . <i>l</i> (x)	:	she loves all things

Propositional logic:

 $\begin{array}{rcl} p & : & it \ is \ raining & \\ s & : & she \ is \ sad & \\ \hline p \supset s & : & if \ it \ is \ raining \ then \ she \ is \ sad \\ p \land s & : & it \ is \ raining \ and \ she \ is \ sad \\ p \lor s & : & it \ is \ raining \ or \ she \ is \ sad \\ \end{array} \right\}^{atomic \ propositions} \left\{ \begin{array}{rcl} \overline{p} & : & it \ is \ not \ raining \\ \hline \overline{s} & : & she \ is \ not \ sad \\ \hline \end{array} \right.$

Modal logic:

¢р	:	it is possible that it is raining
<u>s</u>	:	it is necessary that she is sad

First-order logic:

 $\begin{array}{rcl} l(x) & : & she \ loves \ x \\ h(x,y) & : & x \ hates \ y \end{array} \right\}_{\text{atomic predicates}} \\ \\ \forall x.l(x) & : & she \ loves \ all \ things \\ \exists y. \forall x.h(x,y) \land s & : & there \ exists \ something \ that \ everyone \ hates \ and \ she \ is \ sad \ s$

Observed world





$$\overline{p} : \text{it is not raining} \overset{\textcircled{}}{\odot} \\ s : \text{she is sad} \overset{\textcircled{}}{\odot} \\ V(p) = 0 \qquad V(s) = 1$$

Observed world

$$\overline{p} : \text{ it is not raining } \textcircled{b}{}$$

$$s : \text{ she is sad } \textcircled{b}{}$$

$$V(p) = 0 \qquad V(s) = 1$$

$$V(p \land s) = 0 \qquad V(p \lor s) = 1$$

Observed world

$$\overline{p} : \text{it is not raining} \overset{\textcircled{}_{}_{}}{\overset{}_{}_{}}$$
$$s : \text{she is sad} \overset{\textcircled{}_{}_{}}{\overset{}_{}}$$
$$V(p) = 0 \qquad V(s) = 1$$
$$V(p \land s) = 0 \qquad V(p \lor s) = 1$$
$$V(\Box p) = \qquad V(\diamond s) =$$

Possible world semantics



Possible world semantics

 $w \Vdash \Box A$ iff for all v such that $w \dashrightarrow v, v \Vdash A$



Possible world semantics

 $w \Vdash \Box A$ iff for all v such that $w \dashrightarrow v, v \Vdash A$

 $w \Vdash \Diamond A$ iff there exists v such that $w \dashrightarrow v$ and $v \Vdash A$



How do we reason about such structure?
How do we reason about such structure?

Inference rules:

$$\operatorname{ax}^{n} \frac{1}{a, \overline{a}} = \bigvee_{1}^{n} \frac{A}{A \lor B} = \bigvee_{2}^{n} \frac{B}{A \lor B} = \wedge^{n} \frac{A}{A \land B}$$

How do we reason about such structure?

Inference rules:



How do we reason about such structure?

Inference rules:



Many different approaches

How do we reason about such structure?

Inference rules:



- Many different approaches
- One successful idea!

Syntactical term encoding of the semantical structure

Possible worlds:



Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

0

Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

$$\left[{}^{\mathbf{0}}\bar{p},s,\ldots \right]$$

Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

$$\begin{bmatrix} 0 ar{p}, s, \ldots, \end{bmatrix}^1$$

Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

$$\begin{bmatrix} 0 \overline{p}, s, \ldots, \begin{bmatrix} 1 p, \overline{s}, \ldots \end{bmatrix}$$

Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

$$\begin{bmatrix} 0 \bar{p}, s, \ldots, \begin{bmatrix} 1 p, \bar{s}, \ldots \end{bmatrix}, \begin{bmatrix} 2 & 0 \\ 0 & 0 \end{bmatrix}$$

Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

$$\begin{bmatrix} 0 \overline{p}, s, \ldots, \begin{bmatrix} 1 p, \overline{s}, \ldots \end{bmatrix}, \begin{bmatrix} 2 p, s, \ldots \end{bmatrix}$$

Syntactical term encoding of the semantical structure

Possible worlds:



Nested sequents:

$$\begin{bmatrix} 0 \ \overline{p}, s, \dots, \begin{bmatrix} 1 \ p, \overline{s}, \dots \end{bmatrix}, \begin{bmatrix} 2 \ p, s, \dots, \begin{bmatrix} 3 \ \dots, \begin{bmatrix} 6 \ \dots \end{bmatrix}, \begin{bmatrix} 7 \ \dots \end{bmatrix} \end{bmatrix}, \begin{bmatrix} 4 \ \dots \end{bmatrix}, \begin{bmatrix} 5 \ \dots \end{bmatrix} \end{bmatrix} \end{bmatrix}$$

Nested sequent system:

$ax^n - a, \overline{a}$	$\vee_1^n \frac{A}{A \lor B}$	$\vee_2^n \frac{B}{A \vee B}$	$\wedge^n \frac{A B}{A \wedge B}$

Nested sequent system:

$ax^n - \frac{1}{a, \overline{a}}$	$\vee_1^n \frac{A}{A \lor B}$	$\vee_2^{n} \frac{B}{A \vee B}$	$\wedge^{n} \frac{A - B}{A \wedge B}$
	$\Box^{n} \frac{[{}^{\times}A]}{\Box A}$	$\diamond^{n} \frac{[{}^{y}A,\ldots]}{\diamondA,[{}^{y}\ldots]}$	

Nested sequent system:



An example:

 $\overline{\Diamond a, \Diamond (b \lor c), \Box (\bar{a} \land \bar{b})}$

Nested sequent system:



An example:

 $\Diamond a, \Diamond (b \lor c), \Box (\bar{a} \land \bar{b})$

Nested sequent system:





$$\Box^{\mathsf{n}} \frac{}{\Diamond a, \Diamond (b \lor c), \Box (\bar{a} \land \bar{b})}$$

Nested sequent system:





$$\Box^{\mathsf{n}} \frac{\left[\overline{a} \land \overline{b}\right]}{\Diamond a, \Diamond (b \lor c), \Box (\overline{a} \land \overline{b})}$$

Nested sequent system:



$$\Box^{\mathsf{n}} \frac{\overline{\diamond a, \diamond(b \lor c), [\bar{a} \land \bar{b}]}}{\diamond a, \diamond(b \lor c), \Box(\bar{a} \land \bar{b})}$$

Nested sequent system:



$$\Box^{\mathbf{n}} \frac{\overline{\diamond a, \diamond(b \lor c), [\bar{a} \land \bar{b}]}}{\diamond a, \diamond(b \lor c), \Box(\bar{a} \land \bar{b})}$$

Nested sequent system:



$$\Diamond^n \frac{[A,\ldots]}{\Diamond A,[\ldots]}$$

$$^{\diamond^{n}} \frac{}{ \bigcirc^{a} (b \lor c), [\bar{a} \land \bar{b}]} }{ \diamond^{a} \diamond (b \lor c), \Box (\bar{a} \land \bar{b}) }$$

Nested sequent system:



$$\Diamond^n \frac{[A,\ldots]}{\Diamond A, [\ldots]}$$

$$\Diamond^{\mathsf{n}} \frac{\boxed{[b \lor c]}}{\Diamond a, \Diamond (b \lor c), [\bar{a} \land \bar{b}]} \\ \Diamond a, \Diamond (b \lor c), \Box (\bar{a} \land \bar{b})$$

Nested sequent system:



$$\Diamond^{\mathsf{n}} \frac{\overline{\Diamond a, [b \lor c, \bar{a} \land \bar{b}]}}{\Diamond a, \Diamond (b \lor c), [\bar{a} \land \bar{b}]}$$
$$\Box^{\mathsf{n}} \frac{}{\Diamond a, \Diamond (b \lor c), \Box (\bar{a} \land \bar{b})}$$

Nested sequent system:



$$\Diamond^{\mathbf{n}} \frac{\overline{\diamond \mathbf{a}, \left[b \lor c, \overline{a} \land \overline{b} \right]}}{\diamond a, \diamond (\mathbf{b} \lor \mathbf{c}), \left[\overline{a} \land \overline{b} \right]} \\ \overline{\diamond} a, \diamond (\mathbf{b} \lor \mathbf{c}), \Box (\overline{a} \land \overline{b})$$

Nested sequent system:



$$\Diamond^n \frac{[A,\ldots]}{\Diamond A,[\ldots]}$$

$$\diamondsuit^{n} \frac{\diamondsuit^{n} (b \lor c, \overline{a} \land \overline{b})}{\diamondsuit^{a} (b \lor c), [\overline{a} \land \overline{b}]} \\ \square^{n} \frac{\diamondsuit^{a} (b \lor c), [\overline{a} \land \overline{b}]}{\diamondsuit^{a} (b \lor c), \square (\overline{a} \land \overline{b})}$$

Nested sequent system:



$$\Diamond^n \frac{[A,\ldots]}{\Diamond A,[\ldots]}$$

$$\diamond^{n} \frac{\boxed{[a]}{\diamond a, [b \lor c, \overline{a} \land \overline{b}]}}{\diamond a, \diamond(b \lor c), [\overline{a} \land \overline{b}]}$$
$$\Box^{n} \frac{\diamond a, \diamond(b \lor c), [\overline{a} \land \overline{b}]}{\diamond a, \diamond(b \lor c), \Box(\overline{a} \land \overline{b})}$$

Nested sequent system:



$$\Diamond^{n} \frac{\overline{\left[a, b \lor c, \overline{a} \land \overline{b}\right]}}{\Diamond a, \left[b \lor c, \overline{a} \land \overline{b}\right]}}{ \Diamond a, \Diamond(b \lor c), \left[\overline{a} \land \overline{b}\right]}$$
$$\Box^{n} \frac{\diamond a, \diamond(b \lor c), \left[\overline{a} \land \overline{b}\right]}{\diamond a, \diamond(b \lor c), \Box(\overline{a} \land \overline{b})}$$

Nested sequent system:



$$\Diamond^{\mathsf{n}} \frac{\overline{[a, \mathbf{b} \lor \mathbf{c}, \bar{a} \land \bar{b}]}}{\Diamond a, [b \lor \mathbf{c}, \bar{a} \land \bar{b}]}$$
$$\Box^{\mathsf{n}} \frac{\Diamond a, \Diamond (\mathbf{b} \lor \mathbf{c}), [\bar{a} \land \bar{b}]}{\Diamond a, \Diamond (b \lor \mathbf{c}), \Box (\bar{a} \land \bar{b})}$$

Nested sequent system:



$$\vee_1^n \frac{A}{A \lor B}$$

$$\begin{array}{c} & \bigvee_{1}^{n} \underbrace{[a, b \lor c, \overline{a} \land \overline{b}]}_{\diamond a, [b \lor c, \overline{a} \land \overline{b}]} \\ & \diamond^{n} \underbrace{\frac{\diamond a, [b \lor c, \overline{a} \land \overline{b}]}_{\diamond a, \diamond(b \lor c), [\overline{a} \land \overline{b}]}}_{\diamond a, \diamond(b \lor c), \Box(\overline{a} \land \overline{b})} \end{array}$$

Nested sequent system:





$$\begin{array}{c} & \overline{\bigvee_{1}^{n} \frac{b}{\left[a, b \lor c, \overline{a} \land \overline{b}\right]}} \\ \Leftrightarrow^{n} \frac{}{\diamond a, \left[b \lor c, \overline{a} \land \overline{b}\right]} \\ \xrightarrow{\phi^{n}} \frac{\diamond a, \diamond (b \lor c), \left[\overline{a} \land \overline{b}\right]}{\diamond a, \diamond (b \lor c), \Box (\overline{a} \land \overline{b})} \end{array}$$

Nested sequent system:



$$\square^{n} \frac{\begin{bmatrix} a, b, \overline{a} \land \overline{b} \end{bmatrix}}{\begin{bmatrix} a, b \lor c, \overline{a} \land \overline{b} \end{bmatrix}} \\ \Leftrightarrow^{n} \frac{\begin{bmatrix} a, b \lor c, \overline{a} \land \overline{b} \end{bmatrix}}{\Leftrightarrow a, [b \lor c, \overline{a} \land \overline{b}]} \\ \boxtimes^{n} \frac{\Leftrightarrow a, \diamondsuit (b \lor c), [\overline{a} \land \overline{b}]}{\Leftrightarrow a, \diamondsuit (b \lor c), \square (\overline{a} \land \overline{b})}$$

Nested sequent system:



$$\square^{n} \frac{\begin{bmatrix} a, b, \overline{a} \land \overline{b} \end{bmatrix}}{\begin{bmatrix} a, b \lor c, \overline{a} \land \overline{b} \end{bmatrix}} \\ \diamondsuit^{n} \frac{\begin{bmatrix} a, b \lor c, \overline{a} \land \overline{b} \end{bmatrix}}{\diamondsuit a, [b \lor c, \overline{a} \land \overline{b}]} \\ \square^{n} \frac{\diamondsuit a, \diamondsuit (b \lor c), [\overline{a} \land \overline{b}]}{\diamondsuit a, \diamondsuit (b \lor c), \square (\overline{a} \land \overline{b})}$$

Nested sequent system:



$$\wedge^n \frac{A \quad B}{A \wedge B}$$

$$^{n} \frac{ \bigvee_{1}^{n} \frac{\left[a, b, \overline{a} \land \overline{b}\right]}{\left[a, b \lor c, \overline{a} \land \overline{b}\right]} }{ \diamondsuit^{n} \frac{\left[a, b \lor c, \overline{a} \land \overline{b}\right]}{\diamondsuit a, \left[b \lor c, \overline{a} \land \overline{b}\right]} }{ \diamondsuit^{n} \frac{\diamondsuit^{n} (b \lor c), \left[\overline{a} \land \overline{b}\right]}{\diamondsuit a, \diamondsuit (b \lor c), \left[\overline{a} \land \overline{b}\right]} }$$

Nested sequent system:



$$\wedge^n \frac{A \quad B}{A \wedge B}$$

$$\wedge^{\mathbf{n}} \frac{\overline{\mathbf{a}} \quad \overline{\mathbf{b}}}{\bigvee_{1}^{\mathbf{n}} \frac{[a, b, \overline{\mathbf{a}} \wedge \overline{\mathbf{b}}]}{[a, b \lor c, \overline{a} \wedge \overline{b}]}} \\ \Leftrightarrow^{\mathbf{n}} \frac{\langle \mathbf{a}, [b \lor c, \overline{a} \wedge \overline{b}]}{\diamond \mathbf{a}, [b \lor c, \overline{a} \wedge \overline{b}]} \\ \oplus^{\mathbf{n}} \frac{\langle \mathbf{a}, \phi(\mathbf{b} \lor c), [\overline{a} \land \overline{b}]}{\langle \phi_{a}, \phi(b \lor c), \Box(\overline{\mathbf{a}} \wedge \overline{b})} \end{bmatrix}}$$

Nested sequent system:



$$\wedge^{n} \frac{\overline{[a, b, \overline{a}]} \quad \overline{[a, b, \overline{b}]}}{\bigvee_{1}^{n} \frac{[a, b, \overline{a} \wedge \overline{b}]}{[a, b \vee c, \overline{a} \wedge \overline{b}]}} \\ \Leftrightarrow^{n} \frac{\diamond^{n} \frac{[a, b, \overline{a} \wedge \overline{b}]}{\diamond a, [b \vee c, \overline{a} \wedge \overline{b}]}}{\diamond a, (b \vee c), [\overline{a} \wedge \overline{b}]} \\ \xrightarrow{\circ}^{n} \frac{\diamond (b \vee c), [\overline{a} \wedge \overline{b}]}{\diamond a, (b \vee c), [\overline{a} \wedge \overline{b}]}$$

Nested sequent system:



$$\wedge^{\mathbf{n}} \frac{\overline{[a, b, \overline{a}]} \quad \overline{[a, b, \overline{b}]}}{\bigvee_{1}^{\mathbf{n}} \frac{[a, b, \overline{a} \wedge \overline{b}]}{[a, b \vee c, \overline{a} \wedge \overline{b}]}} \\ \Leftrightarrow^{\mathbf{n}} \frac{\langle \mathbf{a}, [b \vee c, \overline{a} \wedge \overline{b}]}{\langle \mathbf{a}, [b \vee c, \overline{a} \wedge \overline{b}]} \\ \oplus^{\mathbf{n}} \frac{\langle \mathbf{a}, (b \vee c), [\overline{a} \wedge \overline{b}]}{\langle \mathbf{a}, (b \vee c), [\overline{a} \wedge \overline{b}]} \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}, \langle (b \vee c), [\overline{a} \wedge \overline{b}] \\ \hline \rangle \mathbf{a}$$
Nested sequent system:





$$ax^{n} \xrightarrow[]{} \sqrt{[a, b, \overline{a}]} \overline{[a, b, \overline{a}]} \\ \sqrt{[a, b, \overline{a} \land \overline{b}]} } \\ \sqrt{[a, b, \overline{a} \land \overline{b$$

Nested sequent system:



$$ax^{n} \xrightarrow{[a, b, \overline{a}]} \overline{[a, b, \overline{a}]} \xrightarrow{[a, b, \overline{b}]} \\ \wedge^{n} \frac{[a, b, \overline{a} \land \overline{b}]}{\sqrt{n} \frac{[a, b, \overline{a} \land \overline{b}]}{[a, b \lor c, \overline{a} \land \overline{b}]}} \\ \Leftrightarrow^{n} \frac{\Rightarrow^{n} \frac{[a, b, \overline{a} \land \overline{b}]}{\Rightarrow a, [b \lor c, \overline{a} \land \overline{b}]}}{\Rightarrow a, (b \lor c), [\overline{a} \land \overline{b}]} \\ \xrightarrow{[a, b]{}} \frac{\Rightarrow^{n} (b \lor c, \overline{a} \land \overline{b}]}{\Rightarrow a, (b \lor c), \Box(\overline{a} \land \overline{b})}$$

Nested sequent system:





$$ax^{n} \frac{\overline{[a, b, \bar{a}]} ax^{n} \overline{[a, b, \bar{a}]}}{\bigvee_{1}^{n} \frac{[a, b, \bar{a} \wedge \bar{b}]}{[a, b \vee c, \bar{a} \wedge \bar{b}]}} \\ \Leftrightarrow^{n} \frac{\langle a, b, \bar{a} \wedge \bar{b} \rangle}{\langle a, a, b \vee c, \bar{a} \wedge \bar{b} \rangle} \\ \oplus^{n} \frac{\langle a, b \vee c, \bar{a} \wedge \bar{b} \rangle}{\langle a, a, b \vee c, c, \bar{a} \wedge \bar{b} \rangle} \\ ax^{n} \frac{\langle a, b \vee c, \bar{a} \wedge \bar{b} \rangle}{\langle a, a, b \vee c, c, \bar{a} \wedge \bar{b} \rangle}$$

Nested sequent system:

 ${}^{\Gamma}\left\{ \right\} ::= \left\{ \right\} \mid A, {}^{\Gamma}\left\{ \right\} \mid {}^{\Gamma}, [{}^{\Gamma}\left\{ \right\}]$

$ax^n \overline{\Gamma\{a, ar{a}\}}$	$\vee_1^n \frac{\Gamma\{A\}}{\Gamma\{A \lor B\}}$	$\vee_2^n \frac{\Gamma\{B\}}{\Gamma\{A \lor B\}}$	$\wedge^{n} \frac{\Gamma\{A\} \Gamma\{B\}}{\Gamma\{A \land B\}}$
	$\Box^{n} \frac{\Gamma\{[A]\}}{\Gamma\{\Box A\}}$	$\Diamond^{n}\frac{\Gamma_1\{[A,\Gamma_2]\}}{\Gamma_1\{\DiamondA,[\Gamma_2]\}}$	

$$ax^{n} \xrightarrow[]{a, b, \bar{a}]} ax^{n} \overline{[a, b, \bar{a}]} ax^{n} \overline{[a, b, \bar{b}]}$$

$$\wedge^{n} \frac{[a, b, \bar{a} \wedge \bar{b}]}{[a, b \vee c, \bar{a} \wedge \bar{b}]}$$

$$\wedge^{n} \frac{[a, b \vee c, \bar{a} \wedge \bar{b}]}{\Diamond a, [b \vee c, \bar{a} \wedge \bar{b}]}$$

$$\square^{n} \frac{\Diamond a, \Diamond (b \vee c), [\bar{a} \wedge \bar{b}]}{\Diamond a, \Diamond (b \vee c), \square (\bar{a} \wedge \bar{b})}$$

A Unique REarrangement of LIsted ENtities

Choice of focus:

$$\Diamond^{n} \frac{\Diamond a, \left[b \lor c, \bar{a} \land \bar{b}\right]}{\Diamond a, \Diamond(b \lor c), \left[\bar{a} \land \bar{b}\right]} \quad \Diamond^{n} \frac{\Diamond(b \lor c), \left[a, \bar{a} \land \bar{b}\right]}{\Diamond a, \Diamond(b \lor c), \left[\bar{a} \land \bar{b}\right]} \quad \land^{n} \frac{\Diamond a, \Diamond(b \lor c), \left[\bar{a}\right] \quad \Diamond a, \Diamond(b \lor c), \left[\bar{b} \land \bar{b}\right]}{\Diamond a, \Diamond(b \lor c), \left[\bar{a} \land \bar{b}\right]}$$

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Choice of rule:

$$\vee_1^n \frac{A}{A \vee B} \qquad \vee_2^n \frac{B}{A \vee B}$$

Choice of focus:

$$\Diamond^{n} \frac{\Diamond a, \left[b \lor c, \bar{a} \land \bar{b}\right]}{\Diamond a, \Diamond(b \lor c), \left[\bar{a} \land \bar{b}\right]} \quad \Diamond^{n} \frac{\Diamond(b \lor c), \left[a, \bar{a} \land \bar{b}\right]}{\Diamond a, \Diamond(b \lor c), \left[\bar{a} \land \bar{b}\right]} \quad \land^{n} \frac{\Diamond a, \Diamond(b \lor c), \left[\bar{a}\right] \quad \Diamond a, \Diamond(b \lor c), \left[\bar{b}\right]}{\Diamond a, \Diamond(b \lor c), \left[\bar{a} \land \bar{b}\right]}$$

Choice of rule:

$$\vee_1^n \frac{A}{A \vee B} \qquad \vee_2^n \frac{B}{A \vee B}$$

Choice of context:

$$\diamond_1^{\mathsf{n}} \frac{[{}^1A,\ldots], [{}^2\ldots]}{\diamond A, [{}^1\ldots], [{}^2\ldots]} \qquad \diamond_2^{\mathsf{n}} \frac{[{}^1\ldots], [{}^2A,\ldots]}{\diamond A, [{}^1\ldots], [{}^2\ldots]}$$



Proof search space



Proof search space



Proof search space



Focusing provides a way to restrict the proof search space.

Hocus Focus

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Maximal chaining of the decomposition.

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Polarities:	positive formulas	:	P ::=	$a \mid P \lor P \mid \Diamond P \mid \downarrow N$
	negative formulas	:	N ::=	$\bar{a} \mid N \land N \mid \Box N \mid \uparrow P$







Completeness of focusing: If a formula is provable then it has a focused proof.

Control over choices in focused proof is known to improve proof search,



Control over choices in focused proof is known to improve proof search, but also allows for a compact synthetic representation.



$$\vee_{1}^{n} \frac{\bigvee_{2}^{n} \frac{\diamondsuit \frac{[a, \ldots]}{\diamondsuit a, [\ldots]}}{N_{1} \lor \diamondsuit a, [\ldots]}}{(N_{1} \lor \diamondsuit a) \lor N_{2}, [\ldots]}$$









Actual New Utterances (with Passion And Method)

1. Design a system that generates only focused proofs

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 - Reduction of the proof search space

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 - Simple and elegant proof of completeness

Completeness



1. Design a focused labelled framework that can emulate many standard proof systems for modal logic

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- 2. Emulate labelled systems in focused first-order logic (ProofCert)
Contributions:

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comparing formalisms

Contributions:

- 1. Design a focused labelled framework that can emulate many standard proof systems for modal logic
- 2. Emulate labelled systems in focused first-order logic (ProofCert)
 - comparing formalisms
 - checking, (sharing, and reconstructing) proofs

proof in S \longleftrightarrow proof in LMF_{*}

$$\diamond_{\mathsf{Rb}}^{\mathsf{n}} \frac{\Lambda_1\{{}^{\mathsf{x}}[{}^{\mathsf{y}}\Lambda_2], A\}}{\Lambda_1\{{}^{\mathsf{x}}[{}^{\mathsf{y}}\Lambda_2, \Diamond A]\}}$$

proof in S \longleftrightarrow proof in LMF_{*} \longleftrightarrow proof in LKF^a

$$\diamond_{\mathsf{Rb}}^{\mathsf{n}} \frac{\Lambda_1\{{}^{\mathsf{x}}[{}^{\mathsf{y}}\Lambda_2], A\}}{\Lambda_1\{{}^{\mathsf{x}}[{}^{\mathsf{y}}\Lambda_2, \Diamond A]\}}$$

Emulating a nested rule



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Perspectives:

- Applicability and limit of extensions to other modal logics
- Folding: multifocusing and identity of proofs
- Unfolding : Emulation of proof systems for other logics / in different formalisms